

(quasi) differentiation

$$\begin{aligned} d^i(fg) &= (d^i f)g + f(d^i g) \\ &\quad + \sum_{k=1}^i (d^k f)(d^{i-k} g) \end{aligned}$$

§ Introduction

To compute $d^i f$ for $\forall f \in Z(U(M(d, \mathbb{F})))$,
it is desirable that we have $d^i v e^k$
in the form of:

$$d^i v e^k = A(n)_{ij}^i, \text{ where}$$

$$A(n) = \sum_k \underbrace{f_{nik}}_{\substack{\uparrow \\ Z(U(M(d, \mathbb{F})))}} e^k$$

since then

$$\begin{aligned}
& d_j \left((\text{tr } e^m) (\text{tr } e^n) \right) \\
&= (d_j \text{tr } e^m) (\text{tr } e^n) \\
&\quad + (\text{tr } e^m) (d_j \text{tr } e^n) \\
&\quad + \sum_{k=1}^d (d_k^j \text{tr } e^m) (d_k^j \text{tr } e^n) \\
&= A(m)_j (\text{tr } e^n) + (\text{tr } e^m) A(n)_j \\
&\quad + \sum_{k=1}^d A(m)_k^j A(n)_k^j \\
&= \left((\text{tr } e^n) A(m) + (\text{tr } e^m) A(n) + A(m) A(n) \right)_j \\
&= \left(\sum_k (\text{tr } e^n) f_{m,k} e^k + \sum_k (\text{tr } e^m) f_{n,k} e^k \right. \\
&\quad \left. + \sum_{k,l} f_{m,k} f_{n,l} e^{k+l} \right)_j
\end{aligned}$$

§ Recursive formula for $d_j(e^n)^p_q$

on $n = 0, 1, 2, \dots$

$n = 0$

$e^0 = I$ (identity matrix)

$$d_j(e^0)^p_q = d_j I^p_q = 0.$$

$n > 0$

$$d_j(e^n)^p_q = \sum_{v=1}^d d_j(e_v^p(e^{n-1})^v_q)$$

$$= \sum_{v=1}^d (d_j^i e_v^p) (e^{n-1})^v_q$$

$$\uparrow$$
$$d_j^i d_v^p$$

$$+ \sum_{v=1}^d e_v^p d_j^i (e^{n-1})^v_q$$

$$+ \sum_{k=1}^d \sum_{v=1}^d (d_k^i e_v^p) (d_j^k (e^{n-1})^v_q)$$

$$\uparrow$$
$$d_k^i d_j^k$$

$$= (e^{n-1})^i_q d_j^p + \sum_{v=1}^d e_v^p d_j^i (e^{n-1})^v_q$$

$$+ d_j^p (e^{n-1})^i_q$$

Lemma $d_j^i (e^n)_q^p$ has the following form:

$$d_j^i (e^n)_q^p = \sum_{\substack{k_1 \geq 0 \\ k_2 \geq \dots \geq k_n}} \left(a_{k_1, i}^{(n)} (e^k)_q^i (e^e)_j^p \right. \\ \left. + b_{k_1, i}^{(n)} (e^k)_q^p (e^e)_j^i \right),$$

where $a_{k_1, i}^{(n)}$, $b_{k_1, i}^{(n)}$ are

non-negative integers \square

Remark

Lemma $\Rightarrow d_j^i \text{tr } e^n$

$$= \sum_{\substack{k_1 \geq 0 \\ k_2 \geq \dots \geq k_n}} \left(a_{k_1, i}^{(n)} (e^{k_1, i})_q^i \right. \\ \left. + (\text{tr } e^k) b_{k_1, i}^{(n)} (e^e)_j^i \right)$$

$$+ (\text{tr } e^k) b_{k_1, i}^{(n)} (e^e)_j^i$$

\square

Proof of Lemma

Induction on n

$$n=0 \quad d_j^i (e^0)_q^p = 0 \quad \text{clear}$$

Assume that $n > 0$ and

$$d_j^i (e^{n+1})_q^p = \sum_{k \neq i} \left(a_{k i e}^{(n-1)} (e^k)_q^p (e^e)_j^i + b_{k i e}^{(n-1)} (e^k)_q^p (e^e)_j^i \right)$$

Then

$$\begin{aligned} d_j^i (e^n)_q^p &= (e^{n-1})_q^i d_j^p \\ &+ \sum_{r=1}^d e^r d_j^i (e^{n-1})_q^r + d_j^p (e^{n-1})_q^i \\ &= (e^{n-1})_q^i d_j^p + \sum_{k \neq i} a_{k i e}^{(n-1)} \sum_{r=1}^d e^r (e^k)_q^i (e^e)_j^r \\ &\quad + \sum_{k \neq i} b_{k i e}^{(n-1)} (e^{k+1})_q^p (e^e)_j^i \\ &+ \sum_{k \neq i} \left(a_{k i e}^{(n-1)} (e^k)_q^p (e^e)_j^i + b_{k i e}^{(n-1)} (e^k)_q^i (e^e)_j^p \right) \end{aligned}$$

$$\begin{aligned}
& \sum_{v=1}^d e^v (e^k)_q^j (e^e)_j^v \\
&= (e^k)_q^j (e^{e+1})_q^p + \sum_{v=1}^d [e^v_p, (e^k)_q^j] (e^e)_j^v \\
&\quad + \delta_q^p (e^k)_q^j - (e^k)_q^p \delta_q^j \\
&= (e^k)_q^j (e^{e+1})_q^p + \delta_q^p (e^{k+e})_q^j - \underline{(e^k)_q^p (e^e)_j^i}
\end{aligned}$$

$$\begin{aligned}
& \delta_j^i (e^k)_q^p = (e^{k-1})_q^j \delta_j^p \\
&+ \sum_{k \neq e} a_{k \neq e}^{(n-1)} \left((e^k)_q^j (e^{e+1})_q^p + \delta_q^p (e^{k+e})_q^j \right) \\
&+ \sum_{k \neq e} b_{k \neq e}^{(n-1)} \left((e^{k+1})_q^p (e^e)_j^i + (e^k)_q^j (e^e)_j^p \right) \\
&\quad - (*) \quad \square
\end{aligned}$$

The formula (*) gives us recursive determination

$$a_{k \neq e}^{(n-1)}, b_{k \neq e}^{(n-1)} \rightarrow a_{k \neq e}^{(n)}, b_{k \neq e}^{(n)} \rightarrow \dots$$

§ Polynomials associated with $d_j \in e^n$

Since

$$d_j \in e^n = \sum_{\substack{k_1, \dots, k_n \\ k_1 + \dots + k_n = n}} \left(a_{k_1, \dots, k_n}^{(n)} (e^{k_1 e})^j \right. \\ \left. + (1 \in e^k) b_{k_1, \dots, k_n}^{(n)} (e^e)^j \right),$$

we define polynomials $f_k^{(n)}, g_k^{(n)} \in \mathbb{Z}[x]$
($k \in n$) by

$$f_k^{(n)}(x) = \sum_{k_1, \dots, k_n} a_{k_1, \dots, k_n}^{(n)} x^{k_1 + \dots + k_n}$$

$$g_k^{(n)}(x) = \sum_e b_{k_1, \dots, k_n}^{(n)} x^e$$

Then we have

$$d_j \in e^n = f_k^{(n)}(e) \binom{n}{j} + \sum_{k \in n} (1 \in e^k) g_k^{(n)}(e) \binom{n}{j}$$

The formula (*) gives us the recursion

$$\begin{aligned} f^{(n)}(x) &= x^{n-1} + \sum_{k \in \mathbb{Z}} a_{k \in \mathbb{Z}}^{(n-1)} x^{k \in \mathbb{Z} + 1} \\ &\quad + \sum_{k \in \mathbb{Z}} a_{k \in \mathbb{Z}}^{(n-1)} x^{k \in \mathbb{Z}} \\ &= x^{n-1} + x f^{(n-1)}(x) + \sum_{k \in \mathbb{Z}} x^k g_k^{(n-1)}(x) \end{aligned}$$

$$g_0^{(n)}(x) = \sum_{k \in \mathbb{Z}} a_{k \in \mathbb{Z}}^{(n-1)} x^{k \in \mathbb{Z}} = f^{(n-1)}(x)$$

$$g_{k \in \mathbb{Z}}^{(n)}(x) = \sum_{k \in \mathbb{Z}} a_{k \in \mathbb{Z}}^{(n-1)} x^e = g_k^{(n-1)}(x)$$

$$(0 \leq k \in \mathbb{Z})$$

$$\begin{aligned} &= \dots = g_0^{(n-k-1)}(x) \\ &= f^{(n-k-2)}(x) \end{aligned}$$

Therefore, we obtained

$$f^{(n)}(e^x) = \left(f^{(n)}(e) + \sum_{k=0}^{n-1} \binom{n}{k} e^k f^{(n-k-1)}(e) \right) e^x$$

and

$$f^{(n)}(x) = x^{n-1} + x f^{(n-1)}(x) + \sum_{k=0}^{n-2} x^k f^{(n-k-2)}(x)$$

$$f^{(1)}(x), f^{(2)}(x), \dots$$

$$= 0, 1, 2x, 3x^2 + 1, 4x^3 + 4x, \\ 5x^4 + 10x^2 + 4, \dots$$

$$f^{(n)}(x) = \sum_{k: \text{odd}} \binom{n}{k} x^{n-k} \\ = \frac{(x+1)^n - (x-1)^n}{2}$$

Thm

$$j^k \text{tr } e^r = \left(\frac{(e+d)^n - (e-d)^n}{2} \right)$$

$$+ \sum_{k=0}^{n-1} (\text{tr } e^k) \left(\frac{(e+d)^{n-k-1} - (e-d)^{n-k-1}}{2} \right)$$

